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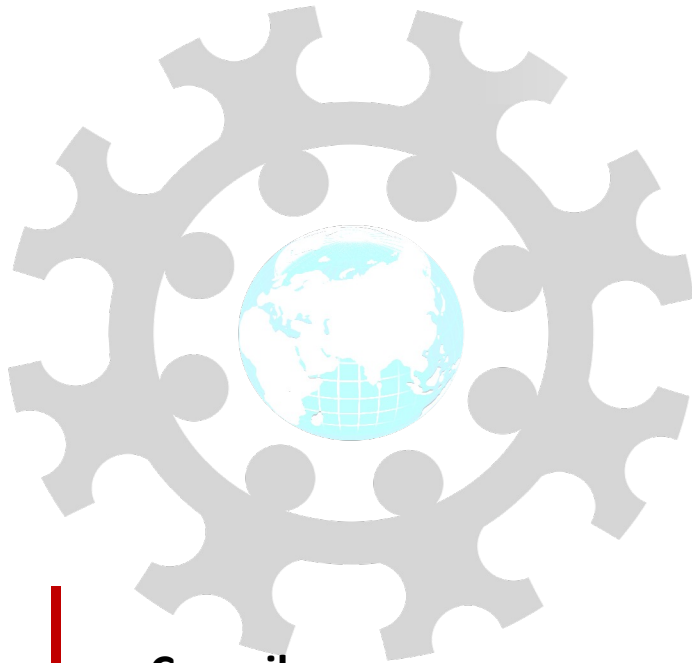
Department of Computer Science

University Institute of Engineering DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

Bachelor of Engineering

Subject Name: System Programming

Subject Code: CST-315



Compilers

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Chapter-1.2

Assembler

- Variants of Assemblers
- Design of two pass assembler



Variants of Assemblers

- There is a list of assemblers ,computer programs that translate assembly language source code into binary programs. Some assemblers are components of a compiler system for a high level language and may have limited or no usable functionality outside of the compiler system.
- - Some assemblers are hosted on the target processor and operating system, while other assemblers (cross-assemblers) may run under an unrelated operating system or processor.

Variants of Assemblers

- **As part of a compiler suite**
- **GNU Assembler (gas):** GPL: many target instruction sets including ARM architecture, Atmel AVR, x86, x86-64, Freescale 68HC11, Freescale v4e, Motorola 680x0, MIPS, PowerPC, IBM System z, TI MSP430.
- **ASxxxx Cross Assembler** (part of the Small Device C Compiler project): GPL: several target instruction sets including Intel 8051, Freescale 68HC08, PIC microcontroller.
- **The Amsterdam Compiler Kit (ACK)** targets many architectures of the 1980s, including 6502, 6800, 680x0, ARM, x86 and Z8000.
- LLVM targets many platforms, however emits no per-target assembly language, instead more high-level typed intermediate representation assembly-like language used.
- Some others self-hosted native-targeted language implementations (like Go, Free Pascal, SBCL) have their own assemblers with multiple targets. They may be used for inline assembly inside language, or even included as a library, but not always suitable for standalone application - no command-line tool exists, or only intermediate representation used as a source, or support for targets very limited.

Single Target assembler

An assembler may have a single target processor or may have options to support multiple processor types. Very simple assemblers may lack features, such as macros, present in more powerful versions.

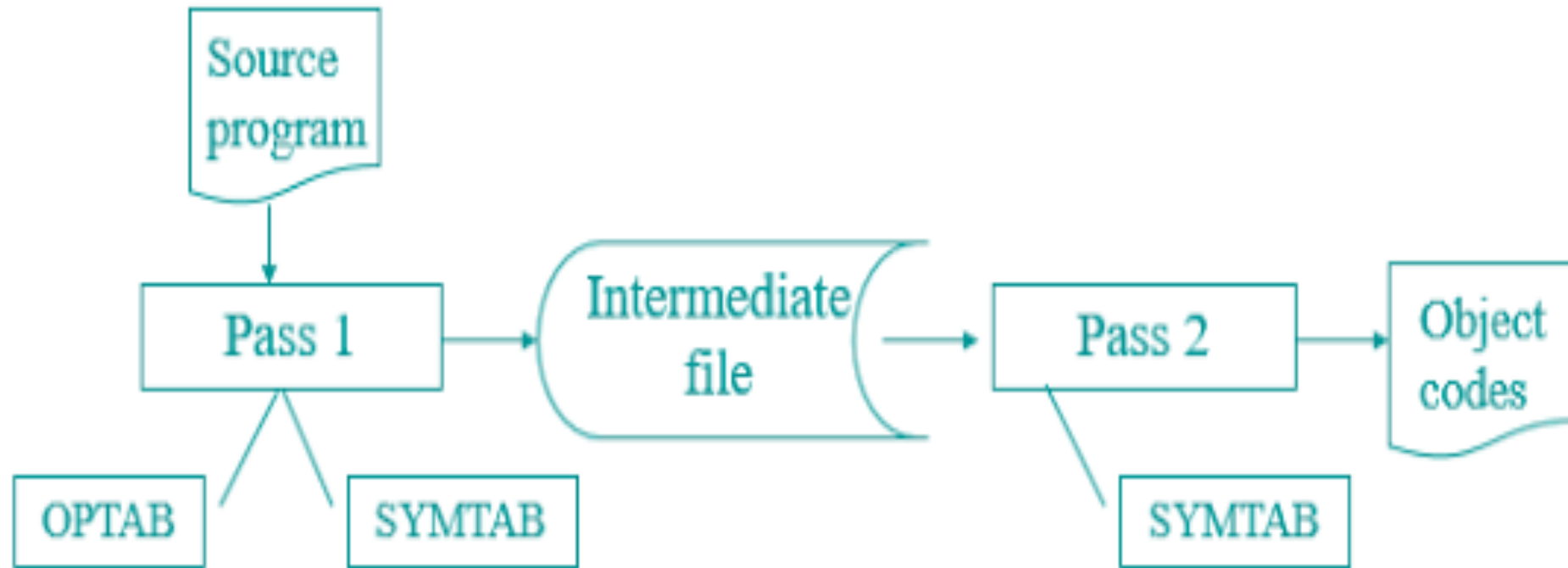
- Various types are:
- 6502 assemblers
- 680x0 assemblers
- ARM assemblers
- Mainframe Assemblers
- POWER, PowerPC, and Power ISA assemblers
- x86 assemblers
- Z80 assemblers
- Other single target assemblers

Design of two pass Assembler

One-pass assembler cannot resolve forward references of data symbols. It requires all data symbols to be defined prior to being used. A two-pass assembler solves this dilemma by devoting one pass to exclusively resolve all (data/label) forward references and then generate object code with no hassles in the next pass. If a data symbol depends on another and this another depends on yet another, the assembler resolved this recursively.

- **Two Pass Assembler**
- *Read from input line*
 - LABEL, OPCODE, OPERAND

Design of 2 pass Assembler



Design of 2 – Pass Assembler

PASS 1:

- Separate the Symbol, Mnemonic opcode, and operand fields
- Build the symbol table
- Perform LC Processing
- Construct Intermediate Representation
- **PASS 2:**
- *SYNTHESIZE THE TARGET PROGRAM*
- Advanced Assembler Directives
- ORIGIN
- EQU

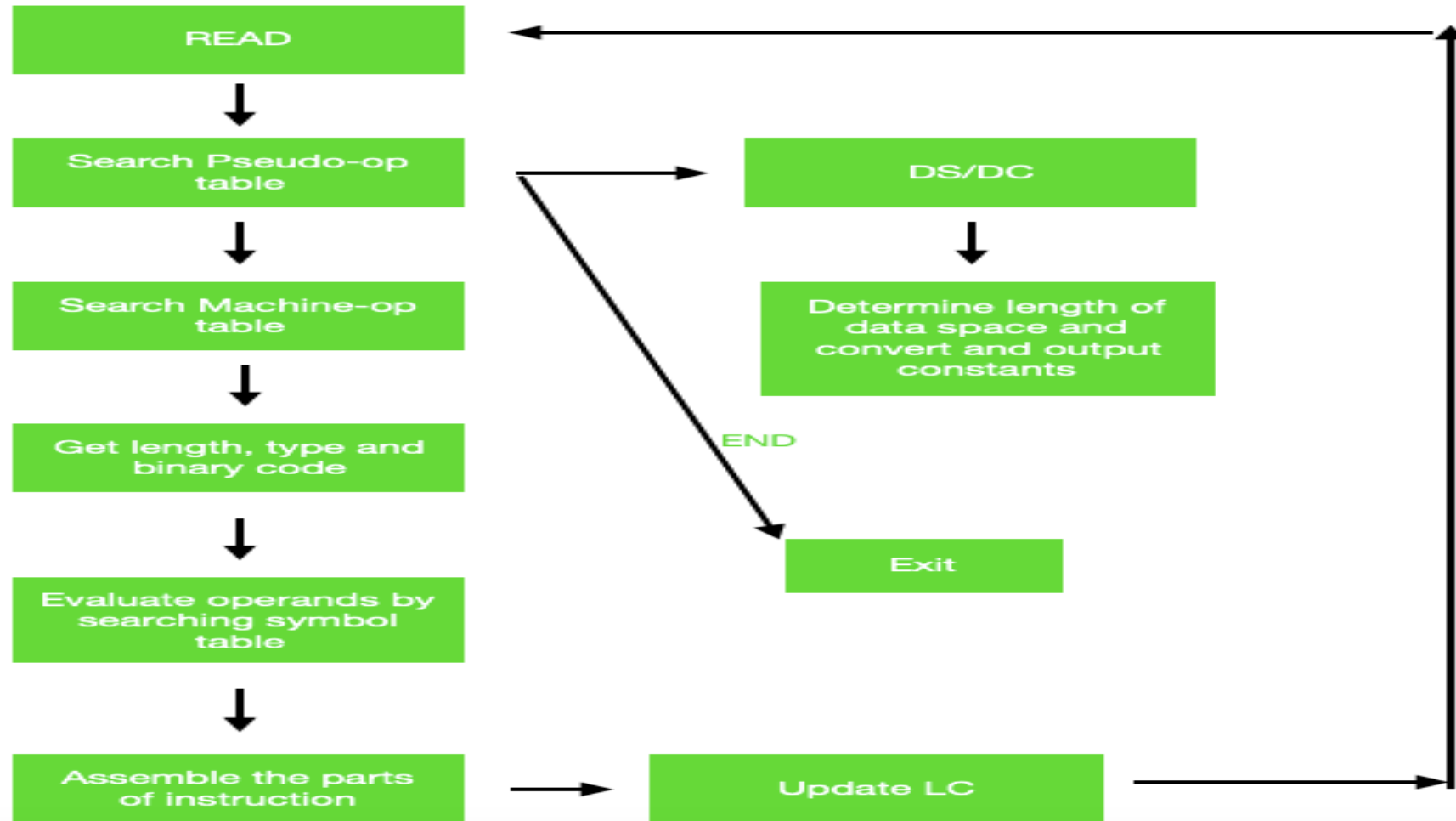
Design of 2 – Pass Assembler

- **EQU**Syntax:
**<Symbol> EQU <Address
Specification>**E.g. MAXLEN EQU 4096 Pass I of Assembler
- Pass I Use following Data Structures
 - OPTAB
 - SYMTAB
 - LITTAB
 - POOLTAB

Design of 2 – Pass Assembler

- 2-pass system is to address the problem of forwarding references — references to variables or subroutines that have not yet been encountered when parsing the source code. A strict 1-pass scanner cannot assemble source code which contains forward references. Pass 1 of the assembler scans the source, determining the size and address of all data and instructions; then pass 2 scans the source again, outputting the binary object code.

Design of 2 – Pass Assembler



Working of Pass-2

- Pass-2 of assembler generates machine code by converting symbolic machine-opcodes into their respective bit configuration(machine understandable form). It stores all machine-opcodes in MOT table (op-code table) with symbolic code, their length and their bit configuration. It will also process pseudo-ops and will store them in POT table(pseudo-op table).
- Various Data bases required by pass-2:
 1. MOT table(machine opcode table)
 2. POT table(pseudo opcode table)
 3. Base table(storing value of base register)
 4. LC (location counter) Take a look at flowchart to understand:

References

- [\[PDF\] Systems Programming and Operating Systems by Dhamdhere - Free Download PDF \(dlscrib.com\)](#)
- [\[PDF\] Principles of Compiler Design By Alfred V. Aho & J.D.Ullman Free Download – Learnengineering.in](#)

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THANK YOU